

# MATH 4LT3/6LT3 Midterm Test Solutions

Midterm Test

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Duration of test: 50 minutes

McMaster University

October 24, 2025

Last Name: \_\_\_\_\_ First Name: \_\_\_\_\_

Student Number: \_\_\_\_\_

This test consists of **four** pages and **four** questions. Please answer all four questions. For all questions, write your answers in the answer booklet that has been provided. Please be sure to include your name and student number on all sheets of paper that you hand in.

**NOTE:** In your solutions you may make use of any theorems or results discussed in the lectures. You may not use other theorems or results, unless you fully justify them. This includes results from the homework assignments, unless otherwise stated.

No aids are allowed.

Each question is worth 5 points; the maximal number of marks is 20.

## Score

Question	1	2	3	4	Total
Score					

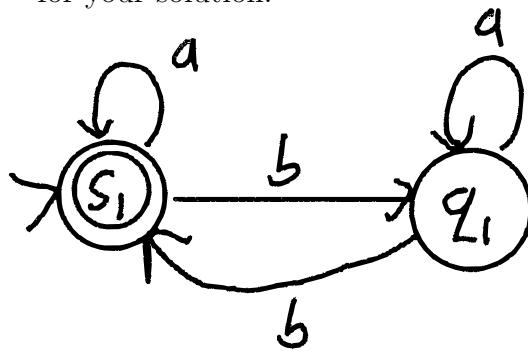
1. Consider the following language over the alphabet  $\Sigma = \{a, b\}$ :

$$L = \{w \in \Sigma^* \mid \#_b(w) \text{ is even or } \#_a(w) = 2\}.$$

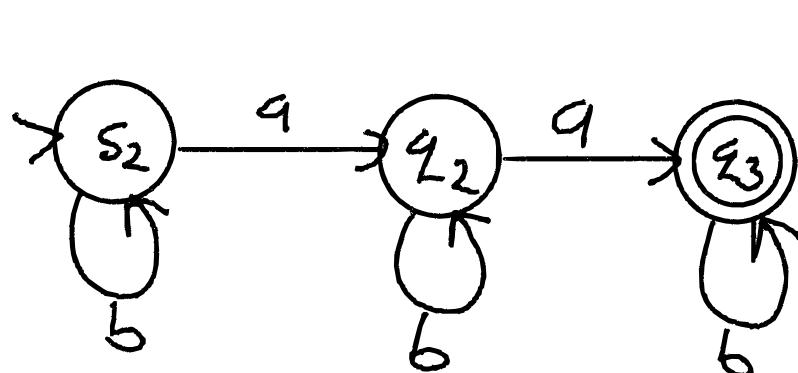
Recall that for a string  $v$  and symbol  $x$ ,  $\#_x(v)$  is the number of occurrences of the symbol  $x$  in the string  $v$ .

So the strings  $aaa$ ,  $bb$ ,  $abbaab$ ,  $babbab$ ,  $abaaba$  all belong to  $L$  and the strings  $abaa$  and  $babb$  do not. The empty string belongs to  $L$ .

Show that  $L$  is a **regular language** by providing an NFA  $M$  with  $L(M) = L$ . Your solution should be in the form of a diagram for the NFA  $M$  that you construct. You do not need to provide a justification for your solution.



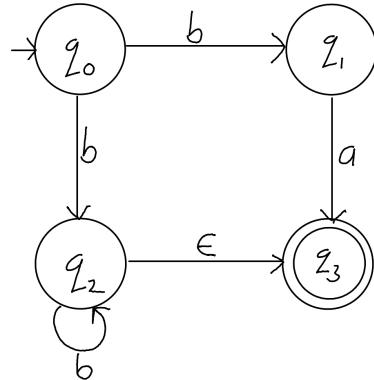
*This part of M  
accepts strings  
that have an odd  
number of b's.*



*this part  
accepts strings  
that have  
exactly 2 a's.*

This NFA has 2 initial states,  $s_1$  and  $s_2$ , and 2 accept states,  $s_1$  and  $q_3$ .

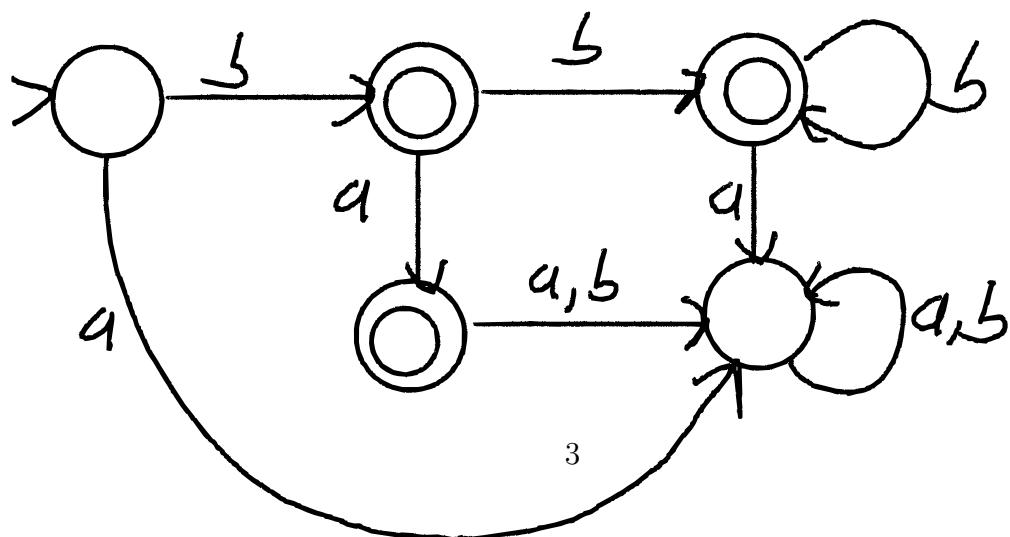
2. Consider the following diagram for the  $\epsilon$ -NFA  $M$  over the alphabet  $\{a, b\}$ :



(a) Describe the language,  $L(M)$ , that this NFA accepts. You do not need to provide a justification for your answer.

$L(M) = \{ba\} \cup \{b^n \mid n \geq 0\}$ .  
 So,  $L(M)$  consists of the string  $ba$  and all strings of  $b$ 's of length at least 1.

(b) Construct a DFA  $M'$  with  $L(M') = L(M)$ . Your solution should be in the form of diagram for the DFA  $M'$ . You do not need to provide a justification for your solution. You do not need to use the subset construction in your solution, but you can if you wish.



For  $1 \leq i < j$ , let  $w = b^i$ .

Then  $a^{3i}w \in L$  and  $a^{3j}w \notin L$ .

3. Consider the following language  $L$  over the alphabet  $\Sigma = \{a, b\}$ :

$$L = \{w \in \Sigma^* \mid \#_a(w) = 3 \times (\#_b(w))\}.$$

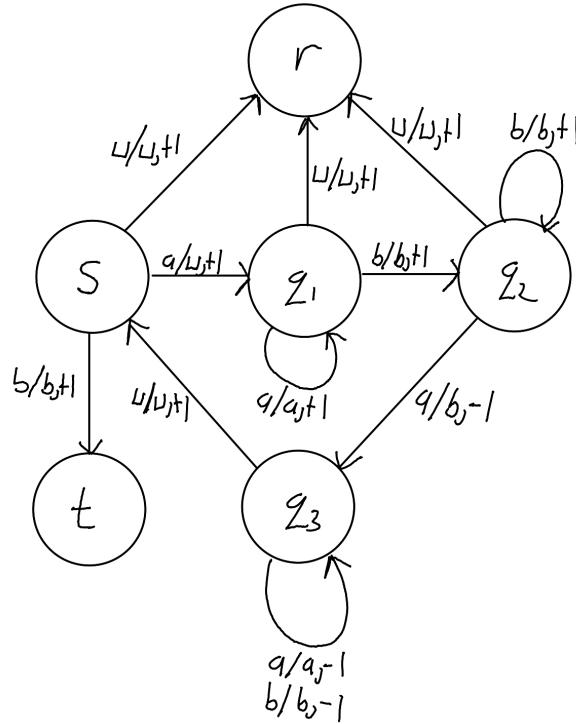
So the strings  $abaa$  and  $baabaaaa$  both belong to  $L$  but the strings  $aaa$ ,  $baab$  do not. The empty string also belongs to  $L$ .

Prove that the language  $L$  is not regular.

Let  $S = \{a^{3n} \mid n \geq 1\}$ . Then  $S$  is an infinite pairwise separable set for  $L$ , and so  $L$  is not regular.

4. The following is a diagram for the Deterministic Turing Machine (DTM)  $T = (Q, \Sigma, \Gamma, s, t, r, \delta)$  where

- $Q = \{s, t, r, q_1, q_2, q_3\}$ ,
- $\Sigma = \{a, b\}$ , and
- $\Gamma = \{a, b, \sqcup\}$ .



$[s, ababb, 1]$

(a) On input  $ababb$ , the DTM  $T$  begins in the initial configuration  $[s, ababb, 1]$ . Write down the next three configurations of the computation of  $T$  on input  $ababb$ .

$[s, ababb, 1] \xrightarrow{T} [q_1, \underline{ababb}, 2] \xrightarrow{T} [q_2, \underline{babbb}, 3]$   
 $\xrightarrow{T} [q_3, \underline{bbb}, 2]$

(b) Which of the following strings are accepted by  $T$ :  $ababb$ ,  $aababb$ ,  $bbbaa$ ? You do not need to provide a justification of your solution.

$ababb$  and  $bbbaa$  are accepted,  
 $aababb$  is rejected.

(c) Give a brief description of  $L(T)$ , the set of strings over  $\Sigma$  that  $T$  accepts. You do not need to provide a justification of your solution.

$L(T)$  consists of all strings  $uv$  that have at least one occurrence of  $b$ , and the number of occurrences of  $a$ 's to the right of the first occurrence of  $b$ 's is at least as big as the number of occurrences of  $a$ 's to the left of the first  $b$ .

$$L(T) = \{ uv \mid \#_a(u) \leq \#_a(v) \}$$